

# Introduction to VC-dimension and its applications in combinatorial problems

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Def.  $(X, \mathcal{F})$  is called a set system if  $X$  is a set and  $\mathcal{F} \subseteq \mathcal{P}(X)$ .

$$\parallel \\ \{Y : Y \subseteq X\}$$

Def. Let  $(X, \mathcal{F})$  be a set system.

(1)  $A \subseteq X$  is shattered by  $\mathcal{F}$  if  $\forall A' \subseteq A, \exists F \in \mathcal{F}$  s.t.  $A \cap F = A'$ .

Equivalently,  $A \subseteq X$  is shattered by  $\mathcal{F}$  if  $\{A \cap F : F \in \mathcal{F}\} = \mathcal{P}(A)$ .

(2)  $\mathcal{F}$  has VC-dimension  $\geq d$  if  $\exists A \subseteq X$  shattered by  $\mathcal{F}$  with  $|A| = d$ .

We write  $VC(\mathcal{F}) \geq d$ .

If  $VC(\mathcal{F}) \geq d$  but not  $VC(\mathcal{F}) \geq d+1$ , then  $VC(\mathcal{F}) = d$ .

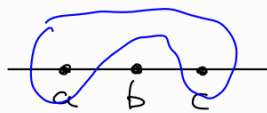
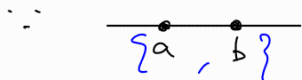
If  $VC(\mathcal{F}) \geq d$  for all  $d \in \mathbb{N}$ , then  $VC(\mathcal{F}) = \infty$ .

(3) For  $A \subseteq X$ ,  $A \cap \mathcal{F} := \{A \cap F : F \in \mathcal{F}\}$ .

Example.

(1) Let  $X = \mathbb{R}$ ,  $\mathcal{F} = \{(-\infty, a) : a \in \mathbb{R}\} \cup \{(b, \infty) : b \in \mathbb{R}\}$ .

Then  $VC(\mathcal{F}) = 2$ .



$\{a, b, c\}$  is not shattered.

( $\{a, b, c\} \cap \mathcal{F} \neq \{a, c\}$ .)

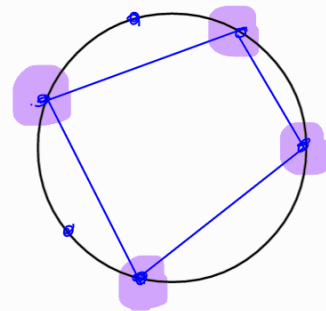
(2) Let  $X = \mathbb{R}^2$ ,  $\mathcal{F} = \{\text{convex polygons}\}$ .

Then  $VC(\mathcal{F}) = \infty$ .

$\therefore$  Let  $A \subseteq \{(x, y) : x^2 + y^2 = 1\}$ ,  $|A| = n$ .

$\forall A' \subseteq A$  with  $|A'| = m \leq n$ ,

$\exists$  convex  $m$ -gon  $G$  s.t.  $A \cap G = A'$ .



Def. Let  $(X, \mathcal{F})$  be a set system.

The shatter function  $\pi_{\mathcal{F}}(n) := \max\{|A \cap \mathcal{F}| : A \subseteq X \text{ and } |A| = n\}$ . ( $\leq 2^n$ )

Lemma. (Sauer-Shelah lemma)

Let  $(X, \mathcal{F})$  be a set system. Assume  $VC(\mathcal{F}) \leq d$ .

Then  $\forall n \geq d$ ,  $\pi_{\mathcal{F}}(n) \leq \sum_{i=0}^d \binom{n}{i}$ . ( $= O(n^d)$ , i.e. polynomial growth.)

Notation.

$$Av(x_1, \dots, x_n; A) := \frac{(1_A(x_1) + \dots + 1_A(x_n))}{n} = \text{relative frequency of } A$$

Theorem. (Weak law of large numbers) (WLLN)

$\forall$  event  $F$ ,  $\forall \varepsilon > 0$ ,  $\forall n > 0$ ,

$$\mathbb{P}\left[|Av(x_1, \dots, x_n; F) - \mathbb{P}[F]| \geq \varepsilon\right] \leq \frac{1}{4n\varepsilon^2} \quad (\rightarrow 0 \text{ as } n \rightarrow \infty).$$

Remark.

$$\mathbb{P}(A \cup B) \leq \mathbb{P}(A) + \mathbb{P}(B)$$

Let  $\mathcal{F}$  be some family of events. By WLLN and union bound,

$$\begin{aligned} & \mathbb{P}\left[\exists \text{ event } F \in \mathcal{F} \text{ s.t. } |Av(x_1, \dots, x_n; F) - \mathbb{P}[F]| \geq \varepsilon\right] \\ &= \mathbb{P}\left[\sup\left\{|Av(x_1, \dots, x_n; F) - \mathbb{P}[F]| : F \in \mathcal{F}\right\} \geq \varepsilon\right] \leq \frac{1}{4n\varepsilon^2} \cdot |\mathcal{F}|. \end{aligned}$$

This "uniform version" w.r.t.  $\mathcal{F}$  has dependency on  $\mathcal{F}$ .

Theorem. (VC-theorem) (Uniform version of the WLLN)

Let  $(X, \mathcal{F})$  be a set system with finite  $X$ .

Regard  $X$  as a probability space with some  $\mathbb{P}$ .  $\forall \varepsilon > 0$ ,

(e.g.  $\mathbb{P}(A) = |A|/|X|$ . Counting probability.)

$$\mathbb{P}\left[\sup\left\{|Av(x_1, \dots, x_n; F) - \mathbb{P}[F]| : F \in \mathcal{F}\right\} > \varepsilon\right] \leq 8\pi_{\mathcal{F}}(n) \exp\left(-\frac{n\varepsilon^2}{32}\right).$$

( $\rightarrow 0$  as  $n \rightarrow \infty$  if  $\pi_{\mathcal{F}}(n)$  has polynomial growth.)

Corollary.

(no dependency on  $\mathcal{F}$ )

$\forall d \in \mathbb{N}$ ,  $\forall \varepsilon > 0$ ,  $\exists N = N(d, \varepsilon)$  s.t.

$\forall$  set system  $(X, \mathcal{F})$  with some  $\mathbb{P}$  on finite  $X$ ,

$VC(\mathcal{F}) \leq d \Rightarrow \exists \varepsilon$ -approximation for  $\mathcal{F}$  of size  $\leq N$ .

$$\left( \exists \{x_1, \dots, x_N\} \subseteq X \text{ s.t. } \forall F \in \mathcal{F}, |Av(x_1, \dots, x_N; F) - \mathbb{P}[F]| < \varepsilon. \right)$$

(possibly with repetitions)

Proof. By VC-theorem,

$$\mathbb{P}\left[\sup\left\{\left|A_N(x_1, \dots, x_N; F) - \mathbb{P}[F]\right| : F \in \mathcal{F}\right\} > \varepsilon\right] \leq 8\pi_{\mathcal{F}}(n) \exp\left(-\frac{n\varepsilon^2}{32}\right).$$

$$\Rightarrow \mathbb{P}\left[\forall F \in \mathcal{F}, \left|A_N(x_1, \dots, x_N; F) - \mathbb{P}[F]\right| \leq \varepsilon\right] \geq 1 - 8\pi_{\mathcal{F}}(n) \exp\left(-\frac{n\varepsilon^2}{32}\right).$$

By SS lemma,  $8\pi_{\mathcal{F}}(n) \exp\left(-\frac{n\varepsilon^2}{32}\right) \rightarrow 0$  as  $n \rightarrow \infty$ .

In particular, with sufficiently large  $N$  depending only on  $d$  and  $\varepsilon$ ,

$$\mathbb{P}\left[\forall F \in \mathcal{F}, \left|A_N(x_1, \dots, x_N; F) - \mathbb{P}[F]\right| < \varepsilon\right] \neq 0.$$

Thus  $\exists$  at least one  $N$ -tuple  $\{x_1, \dots, x_N\}$  s.t.

$$\forall F \in \mathcal{F}, \left|A_N(x_1, \dots, x_N; F) - \mathbb{P}[F]\right| < \varepsilon. \blacksquare$$

## Application to bounding transversal numbers

Def.

Let  $(X, \mathcal{F})$  be a set system.

(1)  $T \subseteq X$  is a transversal of  $\mathcal{F}$  if  $\forall F \in \mathcal{F}, T \cap F \neq \emptyset$ .

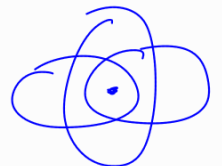
The transversal number of  $\mathcal{F}$ ,

$\tau(\mathcal{F})$  is the smallest size of a transversal.

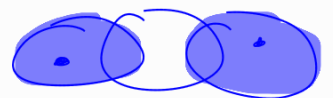
(2)  $\mathcal{G} \subseteq \mathcal{F}$  is a packing if  $\forall F_1 \neq F_2 \in \mathcal{G}, F_1 \cap F_2 = \emptyset$ .

The packing number of  $\mathcal{F}$ ,

$\nu(\mathcal{F})$  is the largest size of a packing.



$$\tau = 2$$



$$\nu = 2$$

Remark.

(1)  $\nu(\mathcal{F}) \leq \tau(\mathcal{F})$ .

( $\therefore$  Let  $\mathcal{G} \subseteq \mathcal{F}$  be a packing with the largest size.

Then  $\forall F \in \mathcal{G}$ , any transversal  $T \subseteq X$  should contain one element from  $F$ .)

(2) It is difficult to calculate the actual transversal/packing number.

(integer programming is NP-hard.)

(3) Very little is known about the reverse direction:

$\tau \leq f(\nu)$  for some fact  $f$  under "reasonable" conditions?

(4) If  $H$  is a hypergraph, then we have a set system

$$X = V(H), \mathcal{F} = E(H) \subseteq \mathcal{P}(V(H)),$$

and a transversal of  $\mathcal{F}$  is called a vertex cover of  $H$ ,  $\tau(H)$ ,

a packing of  $\mathcal{F}$  is called a matching of  $H$ ,  $\nu(H)$ .

(5) If  $H$  is an  $r$ -uniform hypergraph, i.e.  $\forall e \in E(H), |e| = r$ , then

$$\tau(H) \leq r \cdot \nu(H).$$

( $\therefore$  The union of edges from a maximal matching is a vertex cover.)



Ryser's Conjecture (1971)

If  $H$  is an  $r$ -uniform  $r$ -partite hypergraph, then  $\tau(H) \leq (r-1) \cdot \nu(H)$ .

(Solved ~~only~~ for  $r=2$ : König's theorem (1931).)

and  $r=3$  by Ron Aharoni (2001)

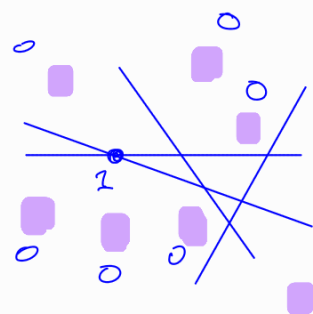
Example.

Let  $X = \mathbb{R}^2$  and  $\mathcal{F}$  be a set of  $n$  lines  $l_1, \dots, l_n$  s.t.

$$\forall i < j < k, |l_i \cap l_j| = 1 \text{ but } l_i \cap l_j \cap l_k = \emptyset.$$

Then  $\nu(\mathcal{F}) = 1$  ( $\therefore$  every two lines intersect),

but  $\tau(\mathcal{F}) \geq \frac{n}{2}$  ( $\therefore$  any point is contained in at most 2 lines).



Def.

(1) Let  $(X, \mathcal{F})$  be a set system with  $X$  finite.

A function  $\phi: X \rightarrow [0, 1]$  is a fractional transversal for  $\mathcal{F}$

$$\text{if } \forall F \in \mathcal{F}, \sum_{x \in F} \phi(x) \geq 1.$$

The size of  $\phi$  is  $\sum_{x \in X} \phi(x)$ .

The fractional transversal number  $\tau^*(\mathcal{F})$  is

the infimum of the sizes of fractional transversals for  $\mathcal{F}$ .

(2) Let  $(X, \mathcal{F})$  be a set system with  $X$  finite.

A function  $\psi: \mathcal{F} \rightarrow [0, 1]$  is a fractional packing for  $\mathcal{F}$

if  $\forall x \in X, \sum_{\{F \in \mathcal{F} : x \in F\}} \psi(F) \leq 1$ .

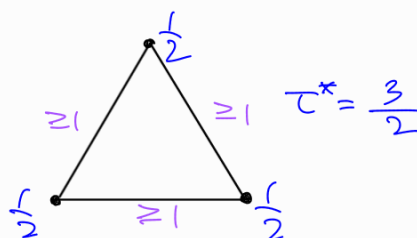
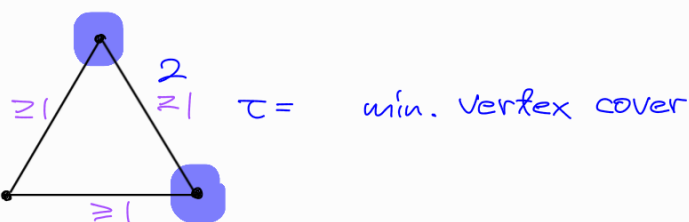
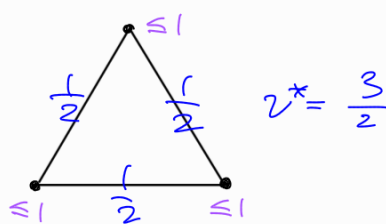
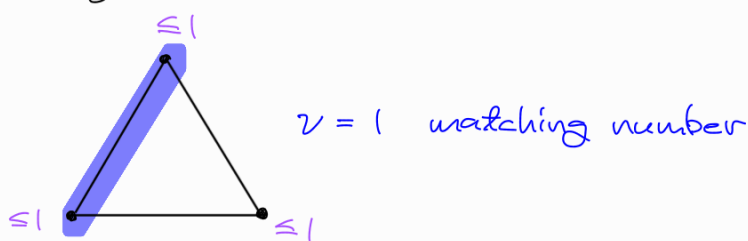
The size of  $\psi$  is  $\sum_{F \in \mathcal{F}} \psi(F)$ .

The fractional packing number  $\nu^*(\mathcal{F})$  is

the supremum of the sizes of fractional packings for  $\mathcal{F}$ .

Example.

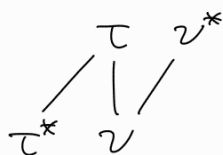
Triangle.  $X = \{a_1, a_2, a_3\}, \mathcal{F} = \{\{a_1, a_2\}, \{a_2, a_3\}, \{a_3, a_1\}\}$ .



Remark.

$$\nu^* \geq \nu \leq \tau \leq \tau^*$$

max     min



Remark.

(1)  $X = [2n], \mathcal{F} = \binom{X}{n}$ . Then  $\tau = n+1$  and  $\tau^* = 2$ .

( $\because$  If  $n$  points  $a_1, \dots, a_n$  are chosen, then  $X \setminus \{a_1, \dots, a_n\} \in \binom{X}{n}$  and

$$\{a_1, \dots, a_n\} \cap (X \setminus \{a_1, \dots, a_n\}) = \emptyset \Rightarrow \tau > n.$$

Assign  $\phi(x) = \frac{1}{n}$  for all  $x \in X \Rightarrow \tau^* \leq 2$ .

$$\sum_{x \in X} \phi(x) = \sum_{x \in \mathcal{F}} \phi(x) + \sum_{x \in X \setminus \mathcal{F}} \phi(x) \geq (n+1) = 2.$$

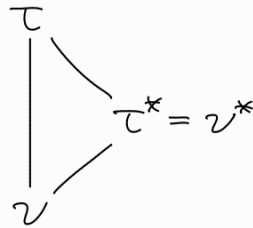
(2)  $X = \binom{[n]}{2}, \mathcal{F} = \left\{ \begin{aligned} &\{\{1, 2\}, \{1, 3\}, \{1, 4\}, \dots\}, \\ &\{\{2, 1\}, \{2, 3\}, \{2, 4\}, \dots\}, \dots \end{aligned} \right\}$

Then  $\nu = 1, \nu^* = \frac{n}{2}$ .

$\mathcal{F}$ : finite

Fact.

(1)  $\mathcal{F}$ : finite  $\Rightarrow \tau^*(\mathcal{F}) = \nu^*(\mathcal{F})$ .



(2) It is easy to calculate or bound the fractional transversal number (linear programming, P.)

Thus it is very useful to bound  $\tau$  in terms of  $\tau^*$ .

Def.

Let  $(X, \mathcal{F})$  be a set system with some  $P$  on  $X$ .

$A \subseteq X$  is called an  $\epsilon$ -net if  $\forall F \in \mathcal{F}$  s.t.  $P[F] \geq \epsilon$ ,  $A \cap F \neq \emptyset$ .

Remark.

Note that every  $\epsilon$ -approximation is an  $\epsilon$ -net.

( $\therefore$  Let  $\{x_1, \dots, x_n\}$  be an  $\epsilon$ -approximation.

$$\forall F \in \mathcal{F}, |A_V(x_1, \dots, x_n; F) - P[F]| < \epsilon$$

$\Rightarrow$  If  $P[F] \geq \epsilon$  but  $\{x_1, \dots, x_n\} \cap F = \emptyset$ , then  $A_V(x_1, \dots, x_n; F) = 0$

and  $|A_V(x_1, \dots, x_n; F) - P[F]| \geq \epsilon$ ,  $\psi$ .)

Fact.

$VC(\mathcal{F}) \leq d \Rightarrow \exists \epsilon$ -net for  $\mathcal{F}$  of size  $\leq Cd \frac{1}{\epsilon} \log \frac{1}{\epsilon}$ .

$\mathcal{F}$ : finite,  $VC(\mathcal{F}) \leq d$

$O(d\tau^* \log \tau^*)$

Proposition.

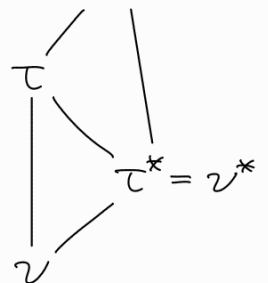
Let  $(X, \mathcal{F})$  be a set system with finite  $X$  and  $VC(\mathcal{F}) \leq d$ .

Then  $\tau(\mathcal{F}) \leq Cd\tau^*(\mathcal{F})\log \tau^*(\mathcal{F})$ , where  $C$ : constant.

Proof.

Let  $\phi: X \rightarrow [0, 1]$  be an optimal fractional transversal, so that

$$\tau^*(\mathcal{F}) = \sum_{x \in X} \phi(x).$$



A function  $P$  defined by  $P(\{x\}) = \frac{\phi(x)}{\tau^*(\mathcal{F})}$  for all  $x \in X$  defines a probability on  $X$ .

Note that  $\forall F \in \mathcal{F}$ ,  $P(F) = \frac{\sum_{x \in F} \phi(x)}{\tau^*(\mathcal{F})} \geq \frac{1}{\tau^*(\mathcal{F})}$ .

Let  $A \subseteq X$  be a  $\frac{1}{\tau^*(\mathcal{F})}$ -net for  $\mathcal{F}$  of size  $\leq C d \tau^*(\mathcal{F}) \log \tau^*(\mathcal{F})$ .

Then  $\forall F \in \mathcal{F}$ ,  $A \cap F \neq \emptyset$ , i.e.,  $A$  is a transversal for  $\mathcal{F}$  of desired size.  $\blacksquare$

Def.

Let  $(X, \mathcal{F})$  be a set system.

$(X^*, \mathcal{F}^*)$  is called the dual set system, where

$X^* = \mathcal{F}$  and  $\mathcal{F}^* = \{\mathcal{F}_x : x \in X\}$  with  $\mathcal{F}_x := \{F \in \mathcal{F} : x \in F\}$ .

The dual VC-dimension of  $\mathcal{F}$ ,  $\text{VC}^*(\mathcal{F})$  is  $\text{VC}(\mathcal{F}^*)$ .

Remark.

(1) Given  $(X, \mathcal{F})$ , the incidence matrix  $M$  is defined as an

$$|X| \times |\mathcal{F}| \text{-matrix s.t. } M_{x,F} = \begin{cases} 1 & \text{if } x \in F, \\ 0 & \text{if } x \notin F. \end{cases}$$

(2) If the incidence matrix of  $(X', \mathcal{F}')$  is the same with  $M$  of  $(X, \mathcal{F})$ ,  
(up to permutations of rows/columns),

then  $(X, \mathcal{F})$  and  $(X', \mathcal{F}')$  are "essentially" the same.

(3) Given  $(X, \mathcal{F})$  with incidence matrix  $M$ ,

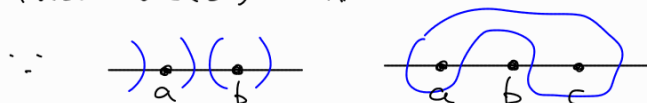
the incidence matrix of  $(X^*, \mathcal{F}^*)$  is  $M^T$ .

( $\because x \in F \iff F \in \mathcal{F}_x$ .)

Example.

Let  $X = \mathbb{R}$ ,  $\mathcal{F} = \{(-\infty, a) : a \in \mathbb{R}\} \cup \{(b, \infty) : b \in \mathbb{R}\}$ .

Then  $\text{VC}(\mathcal{F}) = 2$ .

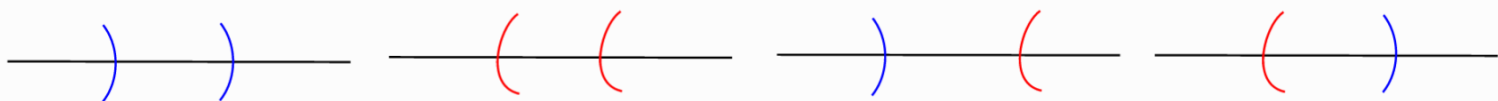


$$X^* = \{(-\infty, a) : a \in \mathbb{R}\} \cup \{(b, \infty) : b \in \mathbb{R}\},$$

$$\mathcal{F}^* = \left\{ \{(-\infty, a) : c \in (-\infty, a)\} \cup \{(b, \infty) : c \in (b, \infty)\} : c \in \mathbb{R} \right\}.$$

$$\ni \mathcal{F}_1 = \left\{ \dots, (-\infty, 10), \dots, (-\infty, 3\sqrt{2}), \dots, (-\infty, 1.01), \dots, \right. \\ \left. \dots, (-10^{10}, \infty), \dots, (0.9999, \infty), \dots \right\}$$

$$VC^*(\mathcal{F}) = 1.$$



### Proposition.

Let  $(X, \mathcal{F})$  be a set system.

Then  $VC^*(\mathcal{F}) < 2^{VC(\mathcal{F})+1}$  and  $VC(\mathcal{F}) < 2^{VC^*(\mathcal{F})+1}$ .

(In particular,  $VC(\mathcal{F}) < \infty$  iff  $VC^*(\mathcal{F}) < \infty$ .)

### Proof.

Assume  $VC(\mathcal{F}) \geq 2^n$ . Then  $\exists A \subseteq X$  shattered by  $\mathcal{F}$  and  $|A| = 2^n$ .

Write  $A = \{a_C : C \subseteq [n]\}$ .

$\forall k \in [n]$ , let  $F_k \in \mathcal{F} = X^*$  s.t.  $A \cap F_k = \{a_C : k \in C \subseteq [n]\}$ .

Then  $\{F_1, \dots, F_n\} \subseteq X^*$  is shattered by  $\mathcal{F}^*$ .

( $\because \forall 1 \leq i_1 < \dots < i_\ell \leq n$ ,  $\{F_1, \dots, F_n\} \cap \mathcal{F}_{a_{\{i_1, \dots, i_\ell\}}} = \{F_{i_1}, \dots, F_{i_\ell}\}$ .)

collection of  $F$ 's s.t.  $F \ni a_{\{i_1, \dots, i_\ell\}}$

### Fact. (Helly's theorem)

Let  $X = \mathbb{R}^d$ ,  $\mathcal{F}$  be some finite family of convex subsets of  $\mathbb{R}^d$ .

Assume that  $\forall \mathcal{F}' \subseteq \mathcal{F}$  s.t.  $|\mathcal{F}'| = d+1$ ,  $\bigcap \mathcal{F}' \neq \emptyset$ .

Then  $\bigcap \mathcal{F} \neq \emptyset$  ( $\Leftrightarrow \tau(\mathcal{F}) = 1$ ).

### Def.

Let  $(X, \mathcal{F})$  be a set system.  $\mathcal{F}$  has fractional Helly number  $k$  if

$\forall \alpha > 0, \exists \beta > 0$  s.t.  $\forall F_1, \dots, F_n \in \mathcal{F}$ ,

if at least  $\alpha$ -fraction of elements  $I \in \binom{[n]}{k}$  satisfy  $\bigcap_{i \in I} F_i \neq \emptyset$ ,

then  $\exists J \subseteq [n]$  s.t.  $|J| \geq \beta n$  and  $\bigcap_{i \in J} F_i \neq \emptyset$ .



### Remark.

(1) Helly's theorem does not hold if convexity is replaced with  $VC(\mathcal{F}) \leq 2$ .

(2) If  $X = \mathbb{R}^d$ ,  $\mathcal{F}$  is some family of convex subsets of  $\mathbb{R}^d$ ,  
then  $\mathcal{F}$  has fractional Helly number  $d+1$  (Katchalski, Liu).

### Theorem (Matousek)

Let  $(X, \mathcal{F})$  be a set system with  $VC^*(\mathcal{F}) \leq k-1$ .

Then  $\mathcal{F}$  has fractional Helly number  $k$ .

### Def.

The VC-density of  $\mathcal{F}$ ,  $vc(\mathcal{F}) = \limsup_{n \rightarrow \infty} \frac{\log(\pi_{\mathcal{F}}(n))}{\log n}$ .

That is, infimum of  $r$ 's s.t.  $\pi_{\mathcal{F}}(n) = O(n^r)$ .

(By SS lemma,  $vc(\mathcal{F}) \leq VC(\mathcal{F})$ .)

VC-codensity  $vc^*(\mathcal{F}) = vc(\mathcal{F}^*)$ .

### Def.

Let  $p \geq q$  be natural numbers and  $(X, \mathcal{F})$  be a set system.

$(X, \mathcal{F})$  satisfies  $(p, q)$ -property if

$\forall \mathcal{F}_0 \subseteq \mathcal{F}$  s.t.  $|\mathcal{F}_0| = p$ ,  $\exists \mathcal{F}_1 \subseteq \mathcal{F}_0$  s.t.  $|\mathcal{F}_1| = q$  and  $\bigcap \mathcal{F}_1 \neq \emptyset$ .

### Remark.

Helly's theorem is equivalent to

$\left( X = \mathbb{R}^d, \mathcal{F}: \text{finite family of convex subsets satisfies } (d+1, d+1)\text{-property} \right)$   
 $\Rightarrow \bigcap \mathcal{F} \neq \emptyset \Leftrightarrow \tau(\mathcal{F}) = 1$ .

### Fact. (Alon, Kleitman)

Let  $p \geq q \geq d+1$  be natural numbers. Then  $\exists N$  s.t.

$\mathcal{F}$ : finite family of convex subsets of  $\mathbb{R}^d$  satisfying  $(p, q)$ -property

$\Rightarrow \tau(\mathcal{F}) \leq N$ .

Theorem. (Alon, Kleitman + Matousek)

Let  $p \geq q \geq d+1$  be natural numbers,  $\mathcal{F}$  satisfy  $(p, q)$ -property.

Then  $\exists N = N(p, q, d)$  s.t.  $(\mathcal{F} : \text{finite and } \text{vc}^*(\mathcal{F}) \leq d \Rightarrow \tau(\mathcal{F}) \leq N)$ .

Main references:

- Artem Chernikov, Model theory and Combinatorics (draft)
- Pierre Simon, A Guide to NIP Theories